# Finding and Understanding Bugs in C Compilers

2021 Most Influential PLDI Paper Award

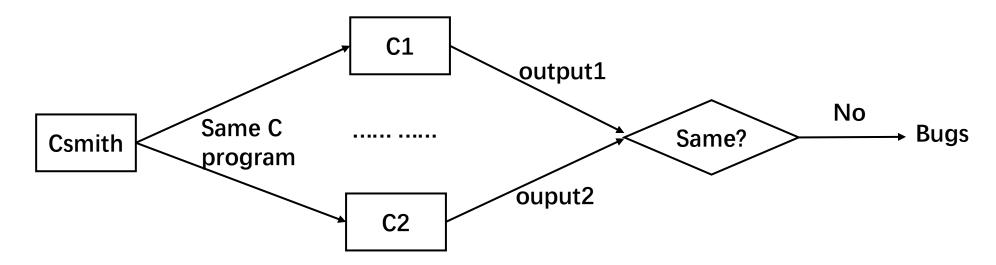
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#### Goal

• Finding bugs in mainstream C compilers like GCC and LLVM.

#### Method

- Randomly generating C programs.
- Differential testing.



#### Contributions

Generating random programs that are expressive.



- Using many C language features.
- Ensuring every program has one single interpretation.

A collection of qualitative and quantitative results about the bugs.

## Randomly Generating Programs

Randomly creating struct type declarations

```
struct X{
   int a;
   float b;
};

struct Y{
   struct X x;
   int c[2];
};
```

```
ExtDef → Specifier FunDec CompSt (recurse)
```

Specifier → TYPE (random) | StructSpecifier

FunDec →ID LP RP (random)

. . . . . .

Exp → ID LP RP (ramdom)

foo is suspended until bar is finished.

#### Randomly Generating Programs

```
struct X{
   int a;
   float b;
};

struct Y{
   struct X x;
   int c[2];
};
```

Output main

Call first generated function

Output checksum

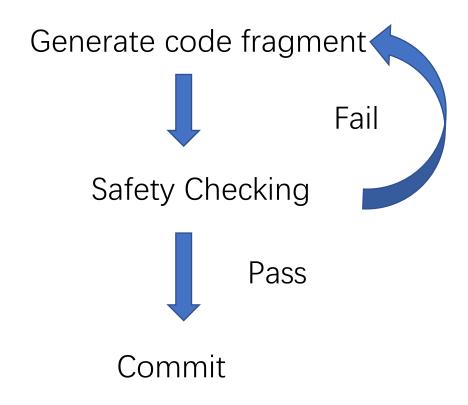
How to ensure one single interpretation?

Static Analysis and Run-time Checking

```
void foo(){
    ....
    bar();
}
```

```
int main(){
    .....
    foo();
    .....
    print(checkSum);
    return 0;
}
```

## Safety Mechanisms



- Integer Safety
- Type Safety
- Pointer Safety
- Effect Safety
- Array Safety
- Initializer Safety

#### Integer safety

Signed Overflow

$$x + 1 > x \rightarrow 1$$
  
INT\_MAX + 1  $\rightarrow$  INT\_MIN

• Shift-past-bitwidth

 $1 \ll 31$  is illegal in C99 with 32-bit ints

Wrapper Functions

# Type safety

Qualifier Safety

```
// object of const-qualified type
const int n = 1;
int* p = (int*)&n;
// undefined behavior
*p = 2;
```

Static Analysis

# Pointer safety

- Null-pointer Dereference.
- Invalid-pointer Dereference.

```
int* p;
int foo(){
    p = 0;
    *p = 1; // null pointer.
    . . . . . .
    int a = 3;
    p = &a;
int bar(){
    int x = *p;//invalid pointer.
```

Pointer Analysis

```
Pts = \{locs, null, invalid\}
No Heap
```

# Effect safety

Unspecified Order

```
func(a(), b());
int a = i++ + ++i;
```

Pointer Analysis

```
Effect = \{Set_{read}, Set_{written}\}
```

Read/Write Conflict between Sequence Points

```
int a = p + func();
```

# Array safety

Indices out of bounds.

For Loop

```
for(int i = 0; i < arr.size(); ++i){
    //not modify i
}</pre>
```

Modulo Operator

```
arr[i % arr.size()];
```

# Initializer safety

Uninitialized Function-scoped Variable

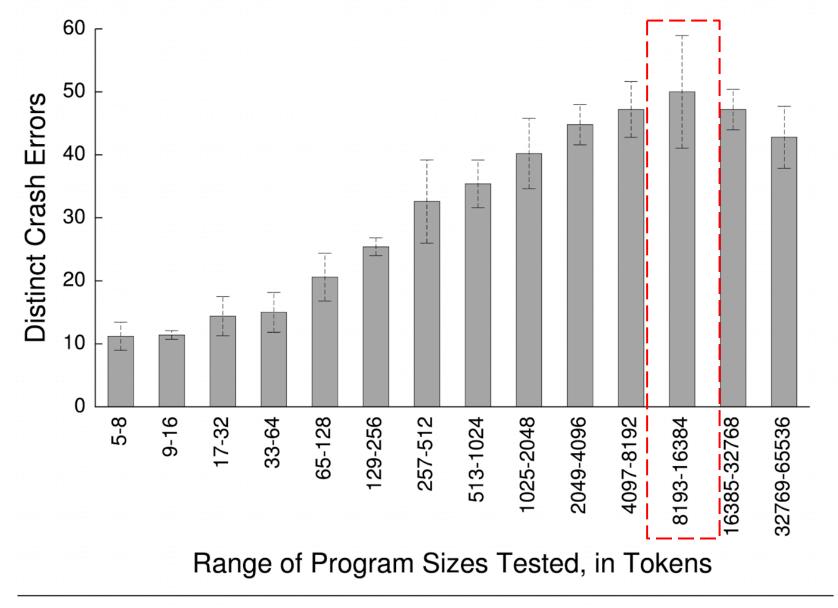
```
int foo(){
    int a;//a is uninitialized
    int x = a + 233;
int foo(){
    int a;
    goto LABEL; //span initializer
    a = 1; //initialized here
    LABEL:
    int x = a + 233;
```

Structurally Ensure Initializing

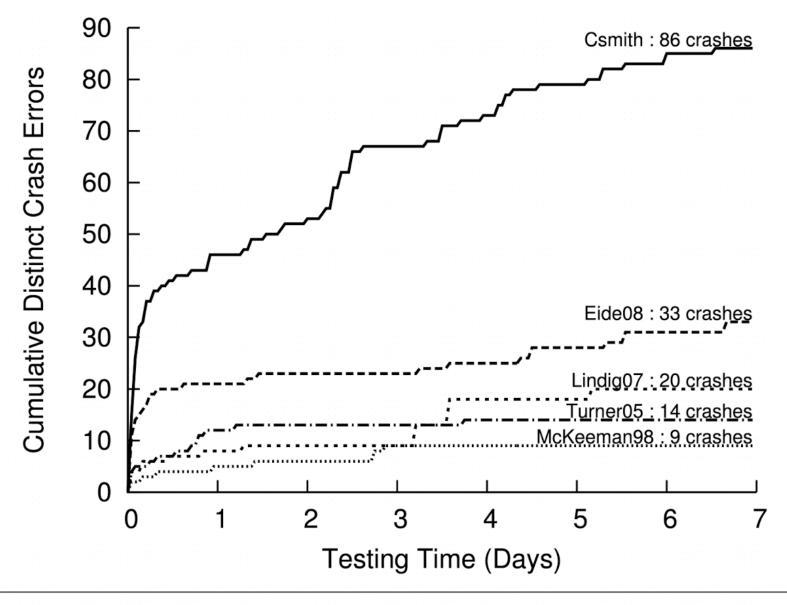
Forbid gotos from spanning initializer

	GCC	LLVM
Crash	2	10
Wrong code	2	9
Total	4	19

**Table 2.** Crash and wrong-code bugs found by Csmith that manifest when compiler optimizations are disabled (i.e., when the -O0 command-line option is used)



**Figure 4.** Number of distinct crash errors found in 24 hours of testing with Csmith-generated programs in a given size range



**Figure 5.** Comparison of the ability of five random program generators to find distinct crash errors

		Line	Function	<b>Branch</b>
		Coverage	Coverage	Coverage
GCC	make check-c	75.13%	82.23%	46.26%
	make check-c & random	75.58%	82.41%	47.11%
	% change	+0.45%	+0.13%	+0.85%
	absolute change	+1,482	+33	+4,471
Clang	make test	74.54%	72.90%	59.22%
	make test & random	74.69%	72.95%	59.48%
	% change	+0.15%	+0.05%	+0.26%
	absolute change	+655	+74	+926

**Table 3.** Augmenting the GCC and LLVM test suites with 10,000 randomly generated programs did not improve code coverage much

Guess: these metrics are too shallow to capture Csmith's effects

# Thank you!